#### Disclosure, Privacy Models, and Privacy Mechanisms

Vicenç Torra

January 2024

Umeå University, Sweden

V. Torra (2022) A guide to data privacy, Springer (Chapter 3)

#### **Outline**

- Disclosure
- Disclosure risk measures
  - Attribute disclosure risk measures
  - Identity disclosure risk measures
- Privacy models

# Introduction to the introduction

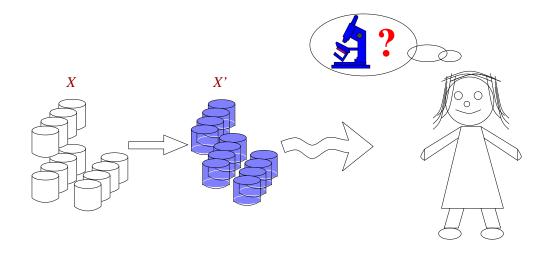
- Disclosure, risk measures, and privacy models
- Protection mechanisms

>

- Data protection mechanisms,
- Privacy-preserving machine learning

Protection mechanisms need to be clearly disassociated of privacy models

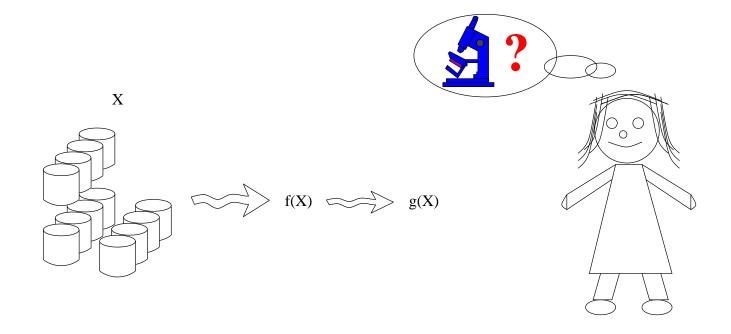
• Privacy for data: data sharing, data publishing



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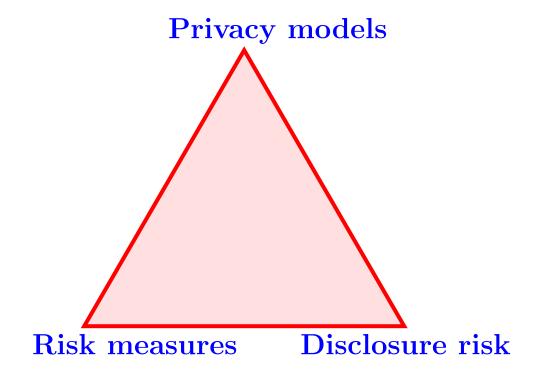
Privacy for computations

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# Introduction

Three strongly related concepts



#### • Def 3.1

- Disclosure takes place when intruders take advantage of the observation and analysis of a release to improve their knowledge on some item of interest.
- Release: data, statistics, data-driven machine learning model

- Disclosure. Attackers take advantage of observations to improve their knowledge on some confidential information about an Iol.
  - $\Rightarrow$  SDC/PPDM: Observe DB,  $\Delta$  knowledge of a particular subject (the respondent in a database)
  - Identity disclosure (entity disclosure). Linkability. Finding Mary in the database.
  - Attribute disclosure. Increase knowledge on Mary's salary.
     also: learning that someone is in the database, although not found.

- Discussion.
  - o Identity disclosure. Avoid.
  - Attribute disclosure. A more complex case. Some attribute disclosure is expected in data mining.

At the other extreme, any improvement in our knowledge about an individual could be considered an intrusion. The latter is particularly likely to cause a problem for data mining, as the goal is to improve our knowledge. (J. Vaidya et al., 2006, p. 7.)

- Another dimension for disclosure.
  - Boolean vs. measurable condition
    - ▷ Boolean: Disclosure either takes place or not. Check, definition holds or not? This includes definitions based on a threshold.
    - Measurable: Disclosure is a matter of degree that can be quantified.
      Some risk is permitted.

- Another dimension for disclosure.
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    - ▷ Boolean: Disclosure either takes place or not. Check, definition holds or not? This includes definitions based on a threshold.
    - ▶ Measurable: Disclosure is a matter of degree that can be quantified.Some risk is permitted.
  - This has implication when selecting a method
    - ▶ Boolean: focus on one performance measure
      - \* minimize information loss (max. utility), minimize execution time
    - $\triangleright$  Measurable: focus on both performance and risk
      - \* multiobjetive optimization problem

• Two dimensions. Privacy models / risk measures

	Attribute disclosure	Identity disclosure
Boolean	Differential privacy Result privacy	k–Anonymity
	Secure multiparty computation	
Quantitative	Interval disclosure	Re-identification (record linkage) Uniqueness

- Attribute disclosure in clusters and cells
  - Aggregates do not avoid disclosure (see motivating example)
  - Inferences depend on the type of summaries indistinguishable/all members of a cell/cluster have the same property
- Types of attacks, for summaries of cells/clusters
  - External attack. An intruder is external to the cell, information is inferred from the analysis of the aggregate.
    - All in a cell have money disorder, so, Mr Scrooge also
    - We learn that Dona Obdúlia was in the psychiatric unit
  - Internal attack. Leakage is caused by intruders (one or coalition)
     that use their own information to learn about the others.
    - All in the cell but one has eating disorders. Mr Scrooge know the disorder of all the others.
    - Dona Obdúlia knows the total income of all other members.

- Attribute disclosure in clusters and cells
  - Anonymization methods based on clustering and aggregation can suffer these types of attacks. E.g., microaggregation, generalization (databases), tabular data protection (aggregates on a few attributes).
  - Privacy models on this type of attacks. Models for tabular data.
     Variants of k-anonymity.

#### Additional discussion.

- Range of opinions. One extreme, privacy is impossible. On the other, risk is overestimated.
- E.g., discussions by de Montjoye et al, replies, and replies to the replies.
- Important point, uniqueness in a sample does not imply uniqueness in the population.

#### Additional discussion.

- Difficulty and success of data privacy depend on the object released (type of data or model) and the level of protection we want to achieve.
- The discussion points out the importance of defining correctly disclosure risk, accurate ways to measure disclosure risk, and welldefined privacy models.
- Measures permit us to evaluate and compare privacy mechanisms (masking methods, privacy-preserving machine learning algorithms as well as different instantiations with different parameters).
- The selection of an appropriate mechanism is key.

- Types of disclosure according to what is disclosed.
  - 1. Attribute disclosure (Section 3.1.3)
  - 2. Identity disclosure (Section 3.1.2)
- Privacy models and measures of disclosure.
  - 1. Measures for attribute disclosure (Section 3.2)
  - 2. Measures for identity disclosure (Section 3.3)
  - 3. Privacy as a measurable condition
    - (a) Uniqueness (Section 3.3.1)
    - (b) Re-identification (Section 3.3.2)
      - i. Data integration: schema and data matching
      - ii. Record linkage algorithms: distance-based and probabilistic RL
      - iii. Generic vs. specific record linkage algorithms
  - 4. Privacy as a Boolean condition
    - (a) k-Anonymity (Section 3.4.2)
    - (b) Differential privacy (Section 3.4.6)
    - (c) Secure multiparty computation (Section 3.4.10)
    - (d) Interval disclosure (Section 3.2.1)

## Risk measures for attribute disclosure

- Difficulties of assessing attribute disclosure
  - We saw:
    - At the other extreme, any improvement in our knowledge about an individual could be considered an intrusion. The latter is particularly likely to cause a problem for data mining, as the goal is to improve our knowledge. (J. Vaidya et al., 2006, p. 7.)
  - For a data-driven model it is difficult to know in what extent we are measuring attribute disclosure and in what extent we are discussing about the quality of the model. The same applies to data releases. Are we producing data of good quality or are we causing attribute disclosure?

- Difficulties of assessing attribute disclosure
  - o Ideally, let M be a data-driven model for attribute A with an excellent generalization capability and no overfitting. Then, the replacement of values A(x) by M(x) does not imply attribute disclosure.

**Algorithm** Rank-based interval disclosure: rid(X, V, V', x, p). **Data**: X: Original file; V: Original attribute; V': Masked attribute; x: record; p: percentage

**Result**: Attribute disclosure for attribute V' of record x

#### begin

```
R(V) := \text{Rank data for attribute } V' i := \text{position of } V'(x) \text{ in } R(V) w := p \cdot |X|/2/100 \text{ (establish the width of the interval)} I(x) = [R[\min(i-w,0)], R[\max(i+w,|X|-1)]] \text{ (interval for } x) rid := V(x) \in I(x) \mathbf{return } rid
```

end

**Data**: X: Original file; V: Original attribute; V': Masked attribute; x: record; p: percentage

**Result**: Attribute disclosure for attribute V' of record x

#### begin

```
sd(V):= standard deviation of V sdid:=|V(x)-V'(x)|\leq p\cdot sd(V)/100 return \underline{sdid}
```

end

## Attribute disclosure for categorical data

- Categorical case: compare V(x) and V'(x)
- Simple attribute disclosure risk (SADR).

$$SADR(X, X') = \frac{|\{x \in X | V(x) = V'(x)\}|}{|X|}.$$

#### Attribute disclosure

- Comparison between V(x) and V'(x) is right?
- If x is changed to x', equality V(x) = V'(x) can be missleading

## Attribute disclosure: model-based (categorical)

**Algorithm 9** Model-based attr disclosure (for x):  $mbd(C_{tr}, A, V, V', x)$ .

**Data**:  $C_{tr}$ : training set; A: algorithm to build a classification model;

V: original attribute; V': masked attribute; x: an example

**Result**: attribute disclosure for x

#### begin

Take the dataset  $C_{tr}$  as input and construct a model of the attribute V using the algorithm A. Let this model be denoted by  $M_{C_{tr},A}$   $V'(x) := M_{C_{tr},A}(x)$ . That is, the application of the classifier  $M_{C_{tr},A}$  to the data x

$$ad := (V'(x) = V(x))$$

return ad

end

## Attribute disclosure: model-based (categorical)

**Algorithm 10** Attribute disclosure for algorithm A: adr(X, X', A).

**Data**: X: original set; X': masked data set; A: for classifiers

**Result**: attribute disclosure risk

#### begin

Define  $C_i^{Tr}$  a partition of X as in cross-validation Compute attribute disclosure risk as follows

$$adr_{A,C,X} := \frac{\sum_{i=1}^{k} |\{x | x \in C_i \text{ and } mbd(C_i^{Tr}, A, V, V', x)\}|}{|X|}.$$

Note that here V(x) refers to the true value of attribute V of x (i.e., according to the original data set)

return  $adr_{A,C,X}$ 

end

## Attribute disclosure: model-based (numerical)

Numerical: Same as for categorical, but with intervals

$$ad := |V(x) - V'(x)| \le p \cdot sd(V)/100$$

## Attribute disclosure: absent attributes (categorical)

**Algorithm 11** Data-driven model-based attr. disclosure ddmbar(X, X', A)

**Data**: X: original set; X': masked data set; A: data mining algorithm

**Result**: attribute disclosure risk in [0,1]

#### begin

Build a data-driven model V from X using algorithm A Build a data-driven model V' from X' using algorithm A Compute attribute disclosure as follows

$$DDMBAR_A = \frac{|\{x|V(x) = V'(x')\}|}{|X|}$$

where x' is the protected version of x return  $DDMBAR_A$ 

#### end

#### Attribute disclosure: discussion

• Boundaries attribute disclosure / data of good quality: difficult

#### • Hints:

- A good estimate requiring a large number of input attributes may be ok, but a good estimate from a few can be problematic. Avoid good inferences from small number of attributes.
- Not all attributes are equally relevant. Confidential and sensitive attributes need to be identified & used for disclosure risk analysis.
- Not all information is always directly available in a database. It may be the case that sensitive information can be inferred from the database but it is not present there. E.g., observance of religious holidays or political orientation may not be an attribute in the database, but the database may contain enough information to estimate these attributes with good accuracy. Attribute disclosure analysis needs to consider whether these cases are relevant.

#### Attribute disclosure: discussion

#### Another perspective:

 In machine and statistical learning, models are expected to generalize data and avoid over-fitting. When a model generalizes correctly and there is no over-fitting, any inference for a particular individual x is due to general properties and not to its particularities. In contrast, bad generalization and over-fitting may imply that inferences are due to memorization and to learning particular features of certain records. When we require good data utility from a machine learning perspective, attribute disclosure should avoid detecting general information found in the data and focus on detecting these particular features of individuals. This has connections with membership attacks, that, in short, try to detect records that are known to have been used in training a model, and they are detected because they are somehow distinguishable from more common ones.

# Attribute disclosure: Through Membership Inference Attacks

- For data-driven models m
- Given x, was x in the training model?
- Is my data used to build m?
- Idea:
  - $\circ$  We build a classifier mia

## Attribute disclosure: Membership Inference Attacks

**Algorithm** Model for membership inference attack:  $mia(C_{tr}, A)$ .

**Data**:  $D^i$ : data sets for building shallow models (each  $D^i$  partitioned into training and testing  $D^{tr}_i$ ,  $D^{te}_i$ ); A: algorithm to build shallow models

Result: Classifier for membership inference attack

```
begin
  sm_i = A(S_i^{tr}) for all i = 1, \ldots, k
 tuples = \emptyset
 for i = 1, \ldots, k do
       forall the x \in D_i^{tr} do
             tuples = \overline{tuples} \cup \{(x, sm_i(x), training)\}
        end
        forall the x \in D_i^{te} do
           tuples = \overline{tuples} \cup \{(x, sm_i(x), no - training)\}
        end
 end
  mia = build-classifier(tuples)
  return mia
```

Vicenç Torra; Privacy models

end

# Risk measures for identity disclosure

- Re-identification. Estimation of correct re-identifications. Theoretically or empirically.
- Uniqueness. Probability that rare combinations in the protected data are also rare in the population.

#### Identity disclosure: Uniqueness

• **Uniqueness.** Risk is defined as the probability that rare combinations of attribute values in the protected data set are indeed rare in the original population.

#### Identity disclosure: Uniqueness

- **Uniqueness.** Risk is defined as the probability that rare combinations of attribute values in the protected data set are indeed rare in the original population.
  - $\circ$  Suitable for sampling ( $\rho(X)$  is a subset of X).
  - For masked data, the same combination will not appear.

#### Identity disclosure: Uniqueness

Measures for identity disclosure: Uniqueness (categorical data/sampling)

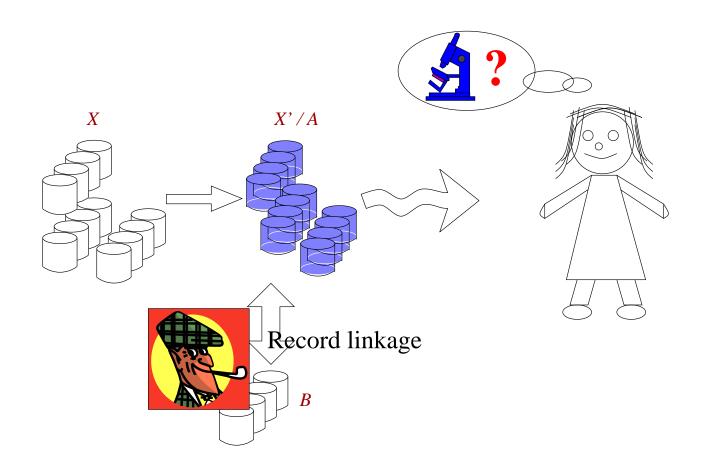
• **File-level uniqueness.** It is defined as the probability that a sample unique (SU) is a population unique (PU). The following expression has been used:

$$P(PU|SU) = \frac{P(PU,SU)}{P(SU)} = \frac{\sum_{j} I(F_j = 1, f_j = 1)}{\sum_{j} I(f_j = 1)}$$

where  $j=1,\ldots,J$  denotes possible values in the sample,  $F_j$  is the number of individuals in the population with key value j (frequency of j in the population),  $f_j$  is the same frequency for the sample and I stands for the cardinality of the selection.

• **Record-level risk uniqueness.** It is defined as the probability that a particular sample record is re-identified (recognized as corresponding to a particular individual in the population).

- Privacy from re-identification. Identity disclosure. Scenario:
  - $\circ$  A: File with the protected data set
  - $\circ$  B: File with the data from the intruder (subset of original X)



**A scenario** for identity disclosure:  $X = id||X_{nc}||X_c|$ 

- Protection of the attributes
  - o **Identifiers.** Usually removed or encrypted.
  - $\circ$  Confidential.  $X_c$  are usually not modified.  $X_c' = X_c$ .
  - $\circ$  Quasi-identifiers. Apply masking method  $\rho$ .  $X'_{nc} = \rho(X_{nc})$ .

- Privacy from re-identification. Identity disclosure.
  - $\circ$  A: File with the protected data set
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How to establish the correct links between the two files? Record linkage algorithms (used in e.g. database integration)

- Privacy from re-identification. Identity disclosure.
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How to establish the correct links between the two files? Record linkage algorithms (used in e.g. database integration)

- Two main types.
  - Distance-based record linkage
  - Probabilistic record linkage

• **Re-identification**. Given  $A = X' = \rho(X)$  and  $B \subset X$ , a measure:

$$Reid(B, A) = \frac{\sum_{b \in B} c(r(b), true(b))}{|B|}.$$
 (1)

#### where

- $\circ true: B \to A$ , for each record b (of the intruder) returns the correct record for re-identification,
- $\circ r: B \to A$ , models the re-identification algorithm. Note: In order to make the definition general, we consider that r returns a probability distribution on A. That is, given a record b in B, it assigns to each record a in A a probability of matching.
- $\circ$  c a function, with c(r(b), true(b)) we evaluate the result for each record in [0,1].

- Re-identification. Examples for Reid(B, A)
  - $\circ$  Algorithm assigns to each record b a single record  $a_b$  in A

$$Reid_d(B, A) = \frac{\sum_{b \in B} |\{b | a_b = true(b)\}|}{|B|}.$$

• Algorithm assigns to each record b an anonymity set in A denoted by  $A_b$ , we model the re-identification algorithm by the distribution  $r(b)[a'] = 1/|A_b|$  for all  $a' \in A_b$ , and r(b)[a'] = 0 for all  $a' \notin A_b$ . Then, we define c(r(b), true(b)) as r(b)[true(b)]. Naturally, we will have that  $c(r(b), true(b)) = r(b)[true(b)] \leq 1/|A_b|$ . Then,

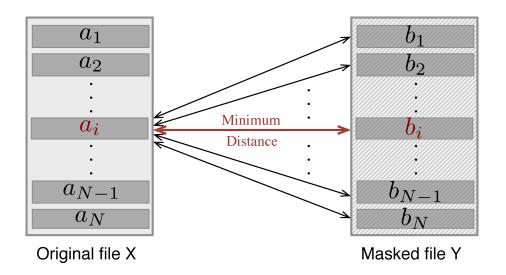
$$Reid_k(B, A) = \frac{\sum_{b \in B} r(b)[true(b)]}{|B|}.$$

Re-identification.
 incorrect links?

Can intruders distinguish correct links from

$$K.Reid(B,A) = \frac{|K|}{|B|} \tag{2}$$

- Distance-based record linkage: d(a,b) with  $a \in A$  and  $b \in B$ .
  - Assign to the record at a minimum distance, ideally an intruder wants for a record i:  $d(a_i,b_j) \geq d(a_i,b_i)$  for all j but due to masking we expect this does not happen



- Algorithm Distance-based record linkage
  - $\circ$  **Input** A: file; B: file **Output** LP: linked pairs; NP: non-linked pairs  $\circ$  For  $a \in A$  $b' = arg \min_{b \in B} d(a, b)$  $LP = LP \cup (a, b')$ for  $b \in B$  such that  $b \neq b'$  $NP := NP \cup (a, b)$ end for o end for  $\circ$  **Return** (LP, NP)

- Probabilistic record linkage: d(a,b) with  $a \in A$  and  $b \in B$ .
  - $\circ$  Classification of pairs of records (a,b) in 3 classes Linked pair, non-linked, clerical pair
  - o How?
    - $\triangleright$  For each pair (a,b), an index is computed using the conditional probabilities
      - $\star P(coincidence|Matching)$ : coincidence between both records when there is matching
      - $\star P(coincidence|Unmatching)$ : coincidence between both records when there is no matching
    - Classification using thresholds

- Probabilistic record linkage: d(a,b) with  $a \in A$  and  $b \in B$ .
  - Computation of

P(coincidence|Matching) and

P(coincidence|Unmatching):

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□ Using EM algorithm

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- Computation of thresholds

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  - Computation of

```
P(coincidence|Matching) and
```

P(coincidence|Unmatching):

- Computation of thresholds
  - ▶ From the probabilities of false positive/negative

```
P(Linkedpair|Unmatching)
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P(Nonlinkedpair|Matching)

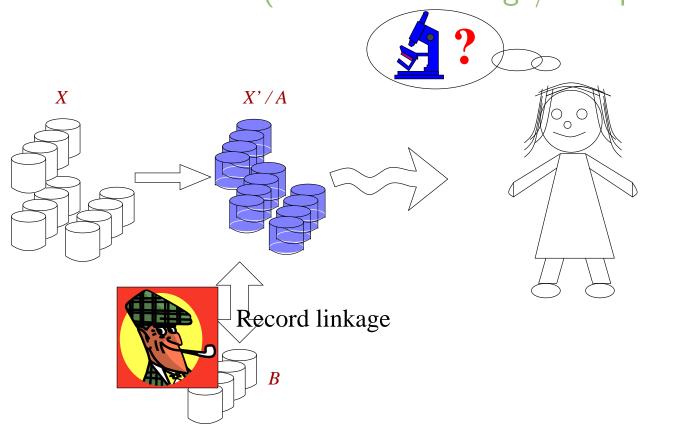
- Flexible scenario for identity disclosure
  - $\circ$  A protected file using a masking method
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    - → intruder with information on only some characteristics

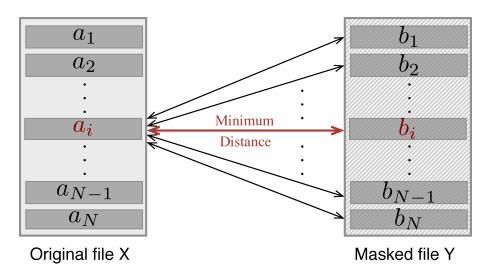
- Flexible scenario for identity disclosure
  - $\circ$  A protected file using a masking method
  - $\circ$  B (intruder's) is a subset of the original file.
    - $\rightarrow$  intruder with information on only some individuals
    - → intruder with information on only some characteristics
  - But also,
    - $\triangleright B$  with a schema different to the one of A (different attributes)
    - Other scenarios. E.g., synthetic data
    - Other type of data: graph data (reidentifying people in a social network)

- Privacy from re-identification. Worst-case scenario (maximum knowledge) to give upper bounds of risk:
  - o transparency attacks (information on how data has been protected)
  - largest data set (original data)
  - best re-identification method (best record linkage/best parameters)



- Privacy from re-identification. Worst-case scenario.
  - $\circ$  ML for distance-based record linkage parameters. (A and B aligned)
    - ▶ Goal: as many correct reidentifications as possible:

for each record 
$$i: d(a_i, b_i) \ge d(a_i, b_i)$$
 for all  $j$ 



 $\blacksquare d(a_i, b_j)$  as average/sum of attribute/variable distances

$$\mathbb{C}_p(diff_1(a_i,b_i),\ldots,diff_n(a_i,b_i))$$

- Privacy from re-identification. Worst-case scenario.
  - $\circ$  ML for distance-based record linkage parameters. (A and B aligned)
    - $\triangleright$  Goal: as many correct reidentifications as possible. But, if error for  $a_i$ :  $K_i = 1$  and  $d(a_i, b_j) + CK_i \ge d(a_i, b_i)$  for all j where d is an aggregated distance  $d(a, b) = \mathbb{C}_p(diff_1, \ldots, diff_n)$ :
    - ▶ Formally,

$$\mathbb{C}_p(diff_1(a_i, b_j), \dots, diff_n(a_i, b_j)) + CK_i \ge \mathbb{C}_p(diff_1(a_i, b_i), \dots, diff_n(a_i, b_i))$$

- Privacy from re-identification. Worst-case scenario.
  - $\circ$  ML for distance-based record linkage parameters. (A and B aligned)
  - Goal: as many correct reidentifications as possible.
  - o Minimize  $K_i$ : minimize the number of records  $a_i$  that fail
- Formalization:

$$Minimize \sum_{i=1}^{N} K_i$$

 $Subject\ to:$ 

$$\mathbb{C}_p(diff_1(a_i, b_j), \dots, diff_n(a_i, b_j)) -$$

$$-\mathbb{C}_p(diff_1(a_i, b_i), \dots, diff_n(a_i, b_i)) + CK_i > 0$$

$$K_i \in \{0, 1\}$$

Additional constraints according to  $\mathbb C$ 

- Privacy from re-identification. Worst-case scenario.
  - $\circ$  ML for distance-based record linkage parameters. (A and B aligned)
  - $\circ$  The case of the weighted mean  $(\mathbb{C}=WM)/\mathsf{W}$ eighted Euclidean
  - Formalization:

$$d^{2}(a,b) = WM_{p}(diff_{1}(a,b), \dots, diff_{n}(a,b))$$

with arbitrary vector 
$$p = (p_1, \dots, p_n)$$
 and  $diff_i(a, b) = ((a_i - \bar{a}_i)/\sigma(a_i) - (b_i - \bar{b}_i)/\sigma(b_i))^2$ 

- Privacy from re-identification. Worst-case scenario.
  - $\circ$  ML for distance-based record linkage parameters. (A and B aligned)
  - $\circ$  The case of the weighted mean ( $\mathbb{C} = WM$ )
  - o Formalization:

$$\begin{aligned} & \textit{Minimize} \ \sum_{i=1}^{N} K_i \\ & \textit{Subject to}: WM_p(\textit{diff}_1(a_i, b_j), \dots, \textit{diff}_n(a_i, b_j)) - \\ & - WM_p(\textit{diff}_1(a_i, b_i), \dots, \textit{diff}_n(a_i, b_i)) + C \ K_i > 0 \\ & K_i \in \{0, 1\} \\ & \sum_{i=1}^{n} p_i = 1 \\ & p_i \geq 0 \end{aligned}$$

- Privacy from re-identification. Worst-case scenario.
  - $\circ$  ML for DBRL parameters: Distances considered  $\mathbb C$ 
    - ▶ Weighted mean.

Weights: importance to the attributes

Parameter: weighting vector n = # attributes

- Privacy from re-identification. Worst-case scenario.
  - $\circ$  ML for DBRL parameters: Distances considered  $\mathbb C$ 
    - ▶ Weighted mean.

Weights: importance to the attributes

Parameter: weighting vector n = # attributes

▷ OWA - linear combination of order statistics (weighted):

Weights: to discard lower or larger distances

Parameter: weighting vector n = # attributes

▶ Bilinear form - generalization of Mahalanobis distance

Weights: interactions between pairs of attributes

Parameter: square matrix:  $n \times n$  (n = # attributes)

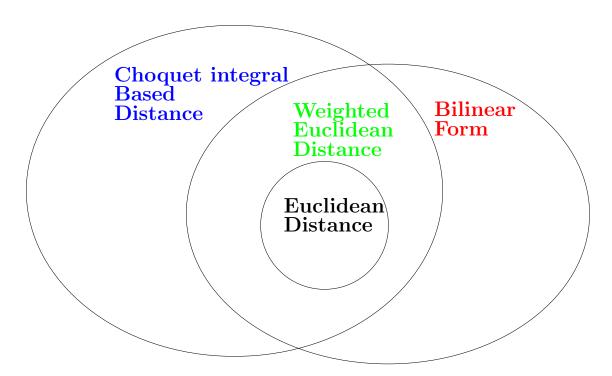
▶ Choquet integral.

Weights: interactions of sets of attributes  $(\mu: 2^X \to [0,1])$ 

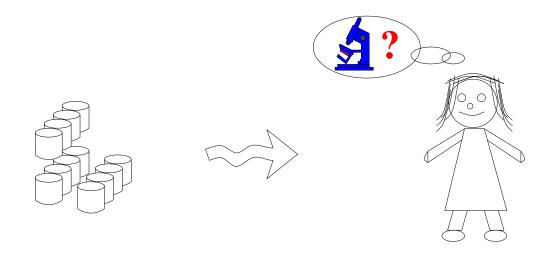
Parameter: non-additive measure:  $2^n - 2$  (n = # attributes)

Distances used in record linkage based on aggregation operators

Graphically



Bilinear form. Quadratic form that generalizes Mahalanobis distance. Choquet integral. A fuzzy integral w.r.t. a fuzzy measure (non-additive measure). Cl generalizes Lebesgue integral. Interactions.

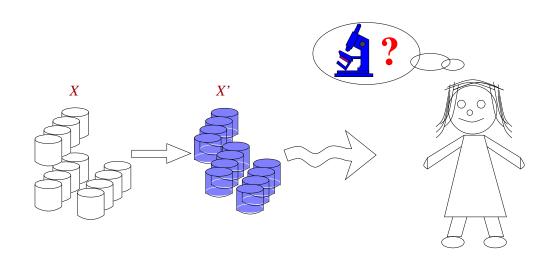


Privacy models. A computational definition for privacy. Examples.

- Reidentification privacy. Avoid finding a record in a database.
- **k-Anonymity.** A record indistinguishable with k-1 other records.
- Secure multiparty computation. Several parties want to compute a function of their databases, but only sharing the result.
- Differential privacy. The output of a query to a database should not depend (much) on whether a record is in the database or not.
- Result privacy. We want to avoid some results when an algorithm is applied to a database.
- Integral privacy. Inference on the databases. E.g., changes have been applied to a database.
- Homomorphic encryption. We want to avoid access to raw data and partial computations.

Privacy models. A computational definition for privacy. Publish a DB

- Reidentification privacy. Avoid finding a record in a database.
- **k-Anonymity.** A record indistinguishable with k-1 other records.
- k-Anonymity, I-diversity. *l* possible categories
- Interval disclosure. The value for an attribute is outside an interval computed from the protected value: values different enough.
- Result privacy. We want to avoid some results when an algorithm is applied to a database.



Privacy models. A computational definition for privacy. Publish a DB

Modify DB X to obtain a DB X' compliant with the privacy model.

Original DB X:

Respondent	City	Age	Illness
DRR	Barcelona	30	Heart attack
ABD	Barcelona	32	Cancer
COL	Barcelona	33	Cancer
GHE	Tarragona	62	AIDS
CIO	Tarragona	65	AIDS
HYU	Tarragona	60	Heart attack

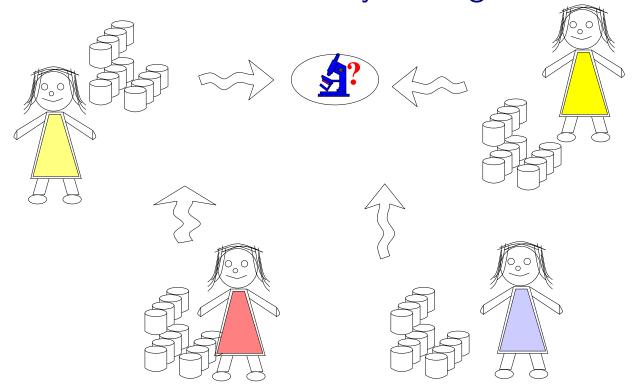
Published DB X':

	City	Age	Illness
	Barcelona	30	Cancer
	Barcelona	30	Cancer
	Barcelona	30	Cancer
	Tarragona	60	AIDS
	Tarragona	60	AIDS
_		_	

- Difficulties
   Naive anonymization does not work, highly identifiable data, high dimensional data
- Examples of successful reidentification attacks
   Sweeney analysis of USA population, data from mobile data, shopping cards, film ratings

Privacy models. A computational definition for privacy. Share a result

• Secure multiparty computation. Several parties want to compute a function of their databases, but only sharing the result.



Privacy models. A computational definition for privacy. Share a result

Compute

$$f(DB_1, DB_2, DB_3, DB_4)$$

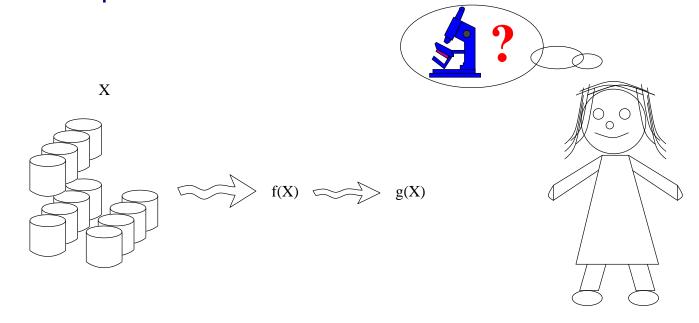
without sharing  $DB_1, DB_2, DB_3, DB_4$ 

• Example: national age mean of hospital-acquired infection patients (hospitals do not want to share the age of their infected patients!)

Difficulties
 Distributed approach (no trusted-third party) – computational cost of solutions

Privacy models. A computational definition for privacy. Compute result

- Differential privacy. The output of a query to a database should not depend (much) on whether a record is in the database or not.
- Integral privacy. Inference on the databases. E.g., changes have been applied to a database.
- Homomorphic encryption. We want to avoid access to raw data and partial computations.



- Difficulties. A simple function can give information on who is in the database
  - E.g., mean salary

# References

#### References

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Vicenç Torra; Privacy models